

Arrays

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Arrays

Definition

An array is a collection of items stored at contiguous memory locations.

The type can be anything: integer, real, character, string or a constructed type.

In this last case, we say that we are dealing with an array of records.

Access to an element : indexing operation

for any i in $[1, n]$: $T(i)$ gives the value of the element having i as index in the array T .

Sorted array:

- set with an order relation.
- for all i in $[1, n-1]$: $T(i) \leq T(i+1)$

Arrays

Static array Vs Dynamic array

Static array

Space allocation is done at the beginning of a treatment.

In technical terms, we say that the space is known at compile time.

Dynamic array

Space allocation is done during execution

The release of the space is then possible during the execution

Arrays

Abstract Machine

ELEMENT (T [i, j, ...]) :

Access to element T[i, j, ...] of array T.

ASS_ELEMENT (T [i, j, ...], Val) :

Assign value Val to element T[i, j, ...].

ALLOC_ARRAY (T) :

Array allocation. The array address is in T.

FREE_ARRAY(T) :

Free the array of address T.

Arrays

Algorithms

4 categories

Traversal : Access by value, Access by position

Update : Insert, Remove

Algorithms on several arrays : Intersection, Union, Merge, Split,...

Sorting arrays

Arrays

Algorithm : Traversal : linear search

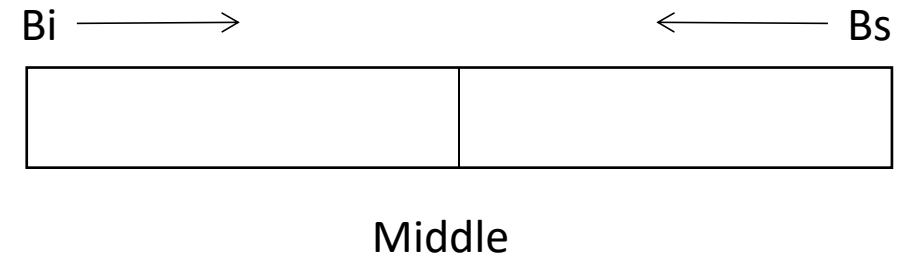
Not sorted array

```
I := 1 ;  
Found := FALSE ;  
WHILE ( I <= N ) AND NOT Found  
  IF ELEMENT ( V [ I ] ) = Val  
    Found := TRUE  
  ELSE  
    I := I + 1  
  ENDIF  
ENDWHILE ;
```

Sorted array

```
I := 1 ;  
Found := FALSE ;  
Stop := FALSE ;  
WHILE ( I <= N ) AND NOT Found AND NOT Stop  
  IF ELEMENT ( V [ I ] ) = Val  
    Found := TRUE  
  ELSE  
    IF ELEMENT ( V [ I ] ) > Val  
      Stop := TRUE  
    ELSE  
      I := I + 1  
    ENDIF  
  ENDIF  
ENDWHILE ;
```

Arrays



Algorithm : Traversal : Binary search

```
Bi := 1 ;  
Bs := N ;  
Found := FALSE ;  
WHILE ( Bi <= Bs ) AND NOT Found  
    Middle := ( Bi + Bs ) / 2 ;  
    IF ELEMENT ( V [ Middle ] ) = Val  
        Found := TRUE  
    ELSE  
        IF Val < ELEMENT ( V [ Middle ] )  
            Bs := Middle - 1  
        ELSE  
            Bi := Middle + 1  
        ENDIF  
    ENDIF  
ENDIF  
ENDWHILE ;
```

Maximal number of
comparisons = $\text{Log}_2(N)$

Arrays

Algorithm : Update

Insert an element after or before a given value

Insert an element at a specific position

Insert an element into a sorted array

Deleting an array element (sorted or not)

Logical delete in a sorted or not sorted array (using a bit lvalid)

Arrays

Maximal number of shifts = N

Algorithm : Inserting a given value into a sorted array T[1..M] holding N elements

```
IF N < M
  {Searching for a position}
  I := 1 ;
  Found := FALSE ;
  WH ( I <= N ) AND NOT Found
    IF ELEMENT ( V [ I ] ) >= Val
      Found := TRUE
    ELSE
      I := I + 1
    ENDIF
  EWH ;

  { Insert  }
  IF NOT Found
    ASS_ELEMENT ( V [ I ] , Val )
  ELSE
    FOR J := N + 1 , I , - 1
      ASS_ELEMENT ( V [ J ] , ELEMENT ( V [ J - 1 ] ) )
    EFOR ;
    ASS_ELEMENT ( V [ I ] , Val )
  ENDIF ;
  N := N + 1
ELSE {N<M}
  WRITE ( 'Overflow' )
ENDIF
```

Arrays

V1 : Array[1..N1] of Integer
V2 : Array[1..N2] of Integer
V3 : ARRAY[1..N1+N2] of Integer

Algorithm : Merging 2 sorted arrays V1[1..N1] and V2[1..N2]

```
I := 1;   J := 1;   K := 0;
WHILE ( I <= N1 ) AND ( J <= N2 )
  K := K + 1;
  IF ELEMENT ( V1 [ I ] ) < ELEMENT ( V2 [ J ] )
    ASS_ELEMENT ( V3 [ K ] , ELEMENT ( V1 [ I ] ) );
    I := I + 1
  ELSE
    ASS_ELEMENT ( V3 [ K ] , ELEMENT ( V2 [ J ] ) );
    J := J + 1
  ENDIF ;
ENDWHILE ;
```

```
WHILE I <= N1
  K := K + 1 ;
  ASS_ELEMENT ( V3 [ K ] , ELEMENT ( V1 [ I ] ) );
  I := I + 1
ENDWHILE ;
WHILE J <= N2
  K := K + 1 ;
  ASS_ELEMENT ( V3 [ K ] , ELEMENT ( V2 [ J ] ) );
  J := J + 1
ENDWHILE
```

Maximal number of
assignments = $N1 + N2$

Arrays

Algorithm : Bubble sort

```
FOR I := 1 , N - 1
  FOR J := N , I + 1 , - 1
    IF ELEMENT ( V [ J - 1 ] ) > ELEMENT ( V [ J ] )
      Temp := ELEMENT ( V [ J - 1 ] );
      ASS_ELEMENT ( V [ J - 1 ] , ELEMENT ( V [ J ] ) );
      ASS_ELEMENT ( V [ J ] , Temp )
    ENDIF
  ENDFOR
ENDFOR
```

Maximal number
of permutations
= $N(N-1)/2$

Arrays

Algorithm : Merge sort

```
Step := 1 ;
WH Step <= N
  FOR I := 1 , N - Step , 2 * Step
    CALL Merge ( I , Step , I + Step , MIN ( Step , N - ( I + Step ) + 1 ) , V2 , N2 ) ;
    CALL Copy ( V2 , N2 , I , I + Step + MIN ( Step , N - ( I + Step ) + 1 ) - 1 ) ;
  EFOR ;
  WRITE ( 'Step=' , Step ) ;
  WRITE ( 'V1=' , V1 ) ;
  Step := 2 * Step
EWH
```

```
V1 V2 V3 V4 V5 V6 V7 V8 V9
[V1] [V2] [V3] [V4] [V5] [V6] [V7] [V8] [V9]

[V1 V2] [V3 V4] [V5 V6] [V7 V8] [V9]

[V1 V2 V3 V4] [V5 V6 V7 V8] [V9]

[V1 V2 V3 V4 V5 V6 V7 V8] [V9]

[V1 V2 V3 V4 V5 V6 V7 V8 V9]
```

Maximal number of
assignments = $N \log_2(N)$

Arrays

Memory representation : n-dimensional array

A typical declaration : $A(a_1:b_1; a_2:b_2; \dots a_n:b_n)$

Order of arrangement of sub-arrays:

$A(a_1, *, *, \dots, *)$, $A(a_1+1, *, *, \dots, *)$, $A(a_1+2, *, *, \dots, *)$,
....., $A(i, *, *, \dots, *)$, $A(b_1, *, *, \dots, *)$

Within each sub-array $A(i, *, *, \dots, *)$, the following order of sub-sub-array is considered:

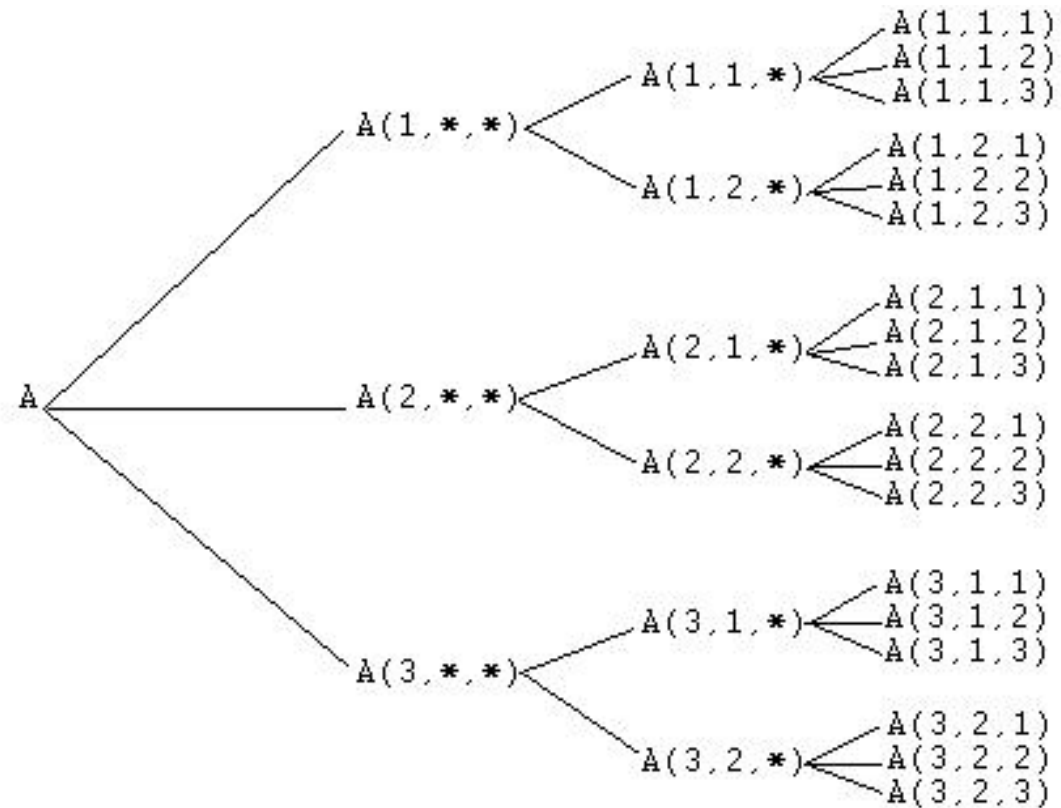
$A(i, a_2, *, *, \dots, *)$, $A(i, a_2+1, *, *, \dots, *)$, $A(i, a_2+2, *, *, \dots, *)$,
....., $A(i, b_2, *, *, \dots, *)$

Etc.

Arrays

Memory representation : n-dimensional array

A(1..3, 1..2, 1..3)



Arrays

Memory representation : n-dimensional array

Address of item $A(i_1, i_2, \dots, i_n)$?

Let $d_i = b_i - a_i + 1$

Address of $A(i_1, *, *, \dots)$: $AD_1 = \text{Base} + (i_1 - a_1) d_2 d_3 \dots d_n$

Address of $A(i_1, i_2, *, \dots)$: $AD_2 = AD_1 + (i_2 - a_2) d_3 d_4 \dots d_n$

Address of $A(i_1, i_2, \dots, i_n)$:

$AD_n = \text{Base} + (i_1 - a_1) d_2 d_3 \dots d_n + (i_2 - a_2) d_3 d_4 \dots d_n + \dots + (i_{n-1} - a_{n-1}) d_n + (i_n - a_n)$

Arrays

Memory representation : n-dimensional array

Address of $A(i_1, i_2, \dots, i_n)$:

$$AD_n = \text{Base} + (i_1 - a_1)d_2d_3\dots d_n + (i_2 - a_2)d_3d_4\dots d_n + \dots + (i_{n-1} - a_{n-1})d_n + (i_n - a_n)$$

Constant part : $\text{Base} - (a_1 \prod_{i=2,n} d_i + a_2 \prod_{i=3,n} d_i + \dots + a_{n-1} d_n + a_n)$

Variable part : $i_1 \prod_{i=2,n} d_i + i_2 \prod_{i=3,n} d_i + \dots + i_{n-1} d_n + i_n$

If $a_1 = a_2 = \dots = a_n = 0$

the address of $A(i_1, i_2, \dots, i_n)$ is

$$i_1d_2d_3\dots d_n + i_2d_3d_4\dots d_n + \dots + i_{n-1}d_n + i_n = \sum_{j=1,n} (i_j \cdot \prod_{i=j+1,n} d_i)$$

Arrays

Implementation (Dynamic / C)

```
#include <stdio.h>
#include <stdlib.h>

/** -Implementation- **\: ARRAY OF INTEGERS**/

/** Arrays **/

typedef int Typeelem_V10i ;
typedef Typeelem_V10i * Typevect_V10i ;

Typeelem_V10i Element_V10i ( Typevect_V10i V , int I1 )
{
    return *(V + (I1-1) );
}

void Ass_element_V10i ( Typevect_V10i V , int I1 ,
Typeelem_V10i Val )
{
    *(V + (I1-1) ) = Val ;
}
```

```
/** Variables of main program **/
Typevect_V10i V;

int main(int argc, char *argv[])
{
    V = malloc(10 * sizeof(int));
    system("PAUSE");
    return 0;
}
```

Arrays

Implementation (Dynamic / PASCAL)

```
PROGRAM My_program;  
  
  { -Implementation- : ARRAY OF INTEGERS }  
  
  { Arrays }  
  TYPE  
    Typeelem_V10I = INTEGER;  
    Typetab_V10I = ARRAY[1..10] OF Typeelem_V10I;  
    Typevect_V10I = ^ Typetab_V10I;  
  
  FUNCTION Element_V10I ( V:Typevect_V10I; I1 : INTEGER ) : Typeelem_V10I ;  
  BEGIN  
    Element_V10I := V^[I1 ];  
  END;  
  
  PROCEDURE Ass_element_V10I ( V :Typevect_V10I; I1 :INTEGER; Val : Typeelem_V10I );  
  BEGIN  
    V^[I1 ] := Val;  
  END;
```

```
{ Declaration part of variables }  
VAR  
  V : Typevect_V10I;  
  
{ Body of main program }  
BEGIN  
  NEW(V);  
  READLN;  
  END.
```